

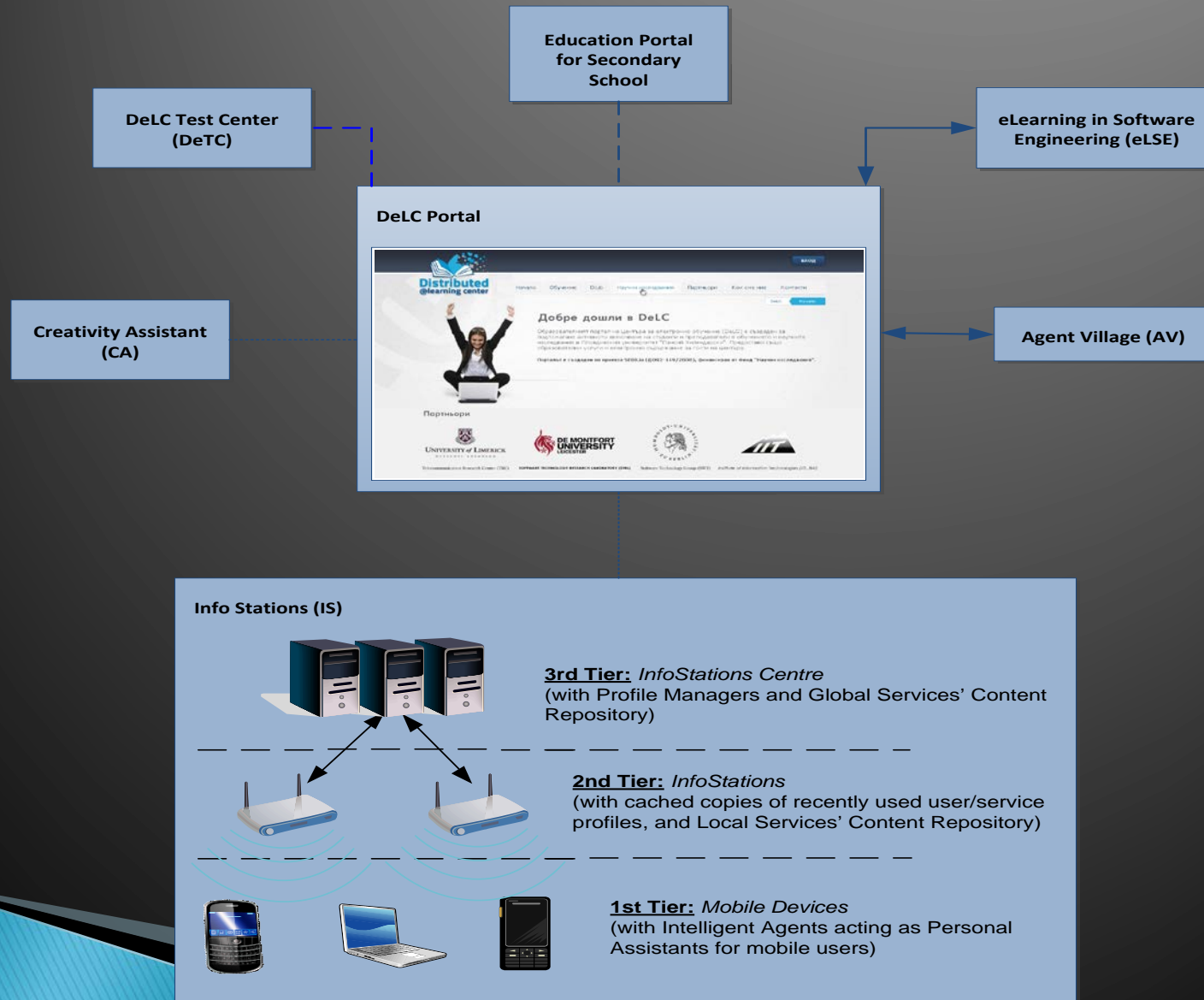
VES – Event management, models and mechanisms (Управление на събития във ВОР – модели и механизми)

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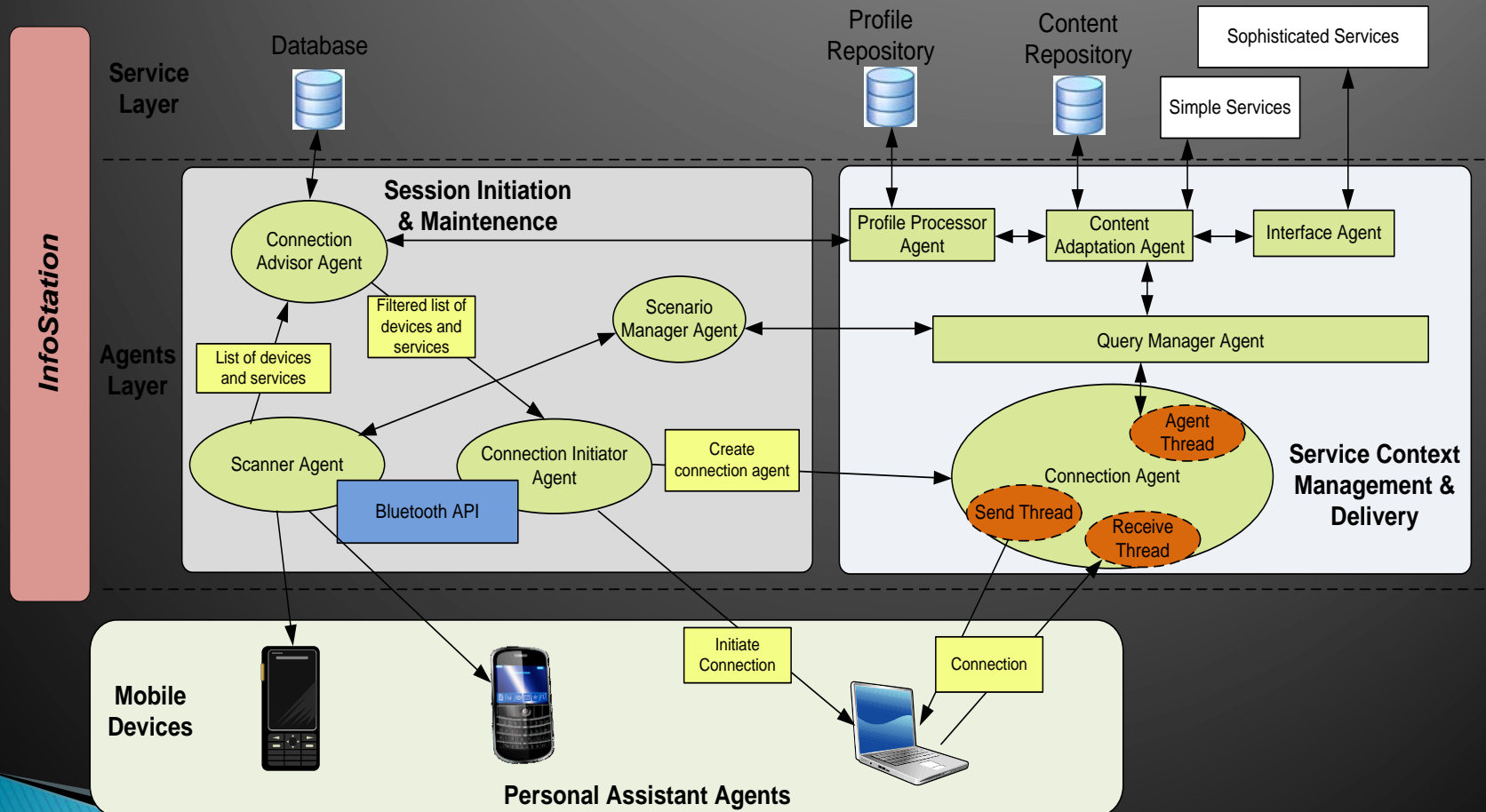
Motivation – DeLC Project

- ▶ Main goal: delivering e-services to support e-learning.
- ▶ DeLC architecture: distributed system, containing fixed and mobile nodes.
- ▶ Mobile node: providing mobile access to the services, through intelligent wireless network based on InfoStation architecture.

Infrastructure of DeLC



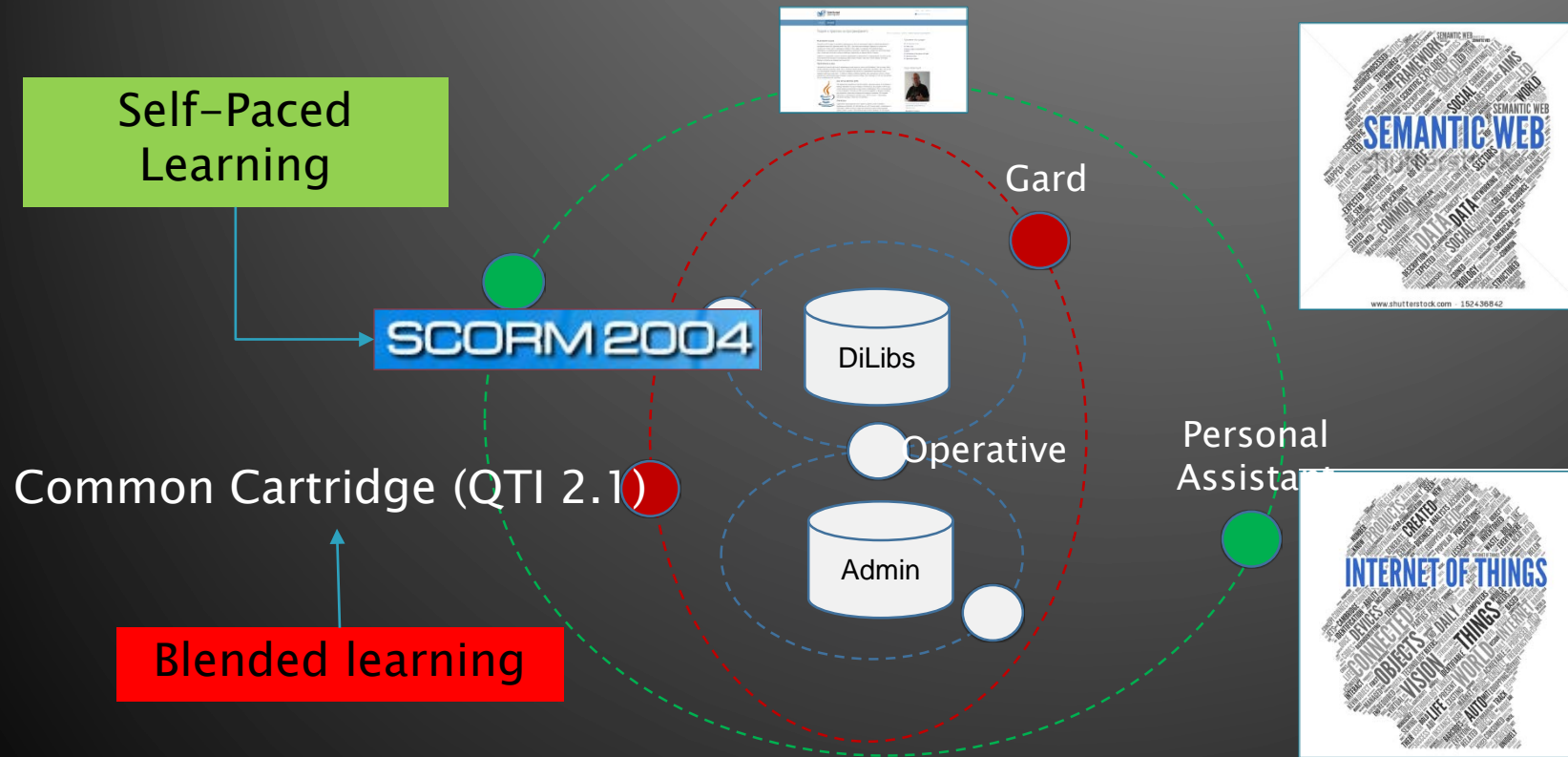
Mobile node middleware



From DeLC to VES

- ▶ New tendency in the development of Internet and Web:
 - Internet of Things – stimulating the origin of cyber-physical systems which will lead to essential consequences in the following years.
 - Semantic web.
- ▶ Develop VES with the following features:
 - Intelligent.
 - Context-aware.
 - Scenario-oriented.
 - Controlled infrastructure.
- ▶ Lifelong learning support

Virtual Education Space



VES in more details: Building a Virtual Education Space, WMSCI 2015, July 12 – 15 – Orlando, Florida, USA

Virtual eLearning Space specifics

- ▶ Functionality has to be supported by agents.
- ▶ VES has a context-aware architecture (adaptation, personalization).
- ▶ Functionality of VES is a non constant set of resources.
- ▶ There is a basic functionalities without with VES cannot exists.
- ▶ Specific components in VES have to react to the changes of the environment.

Main goal

- ▶ To expand the middleware with intelligent agents, which are able to detect and manage time aspects of delivering educational services and content in distributed InfoStation network.
- ▶ Preparing DeLC for being part of Virtual eLearning Space (VES).

Time aspects

- ▶ Scenarios
 - They specify the functionality of our communication environment (InfoStations)
- ▶ During the execution of a service different local events could happened :
 - Getting in/out of range of an IS.
 - Change communication protocol.
 - Change the mobile device.
- ▶ The existing middleware could react to various events, but it is unable to represent them in time order.
 - No management mechanism.
- ▶ In the scenario point of view the problem is to manage scenario change and execution.

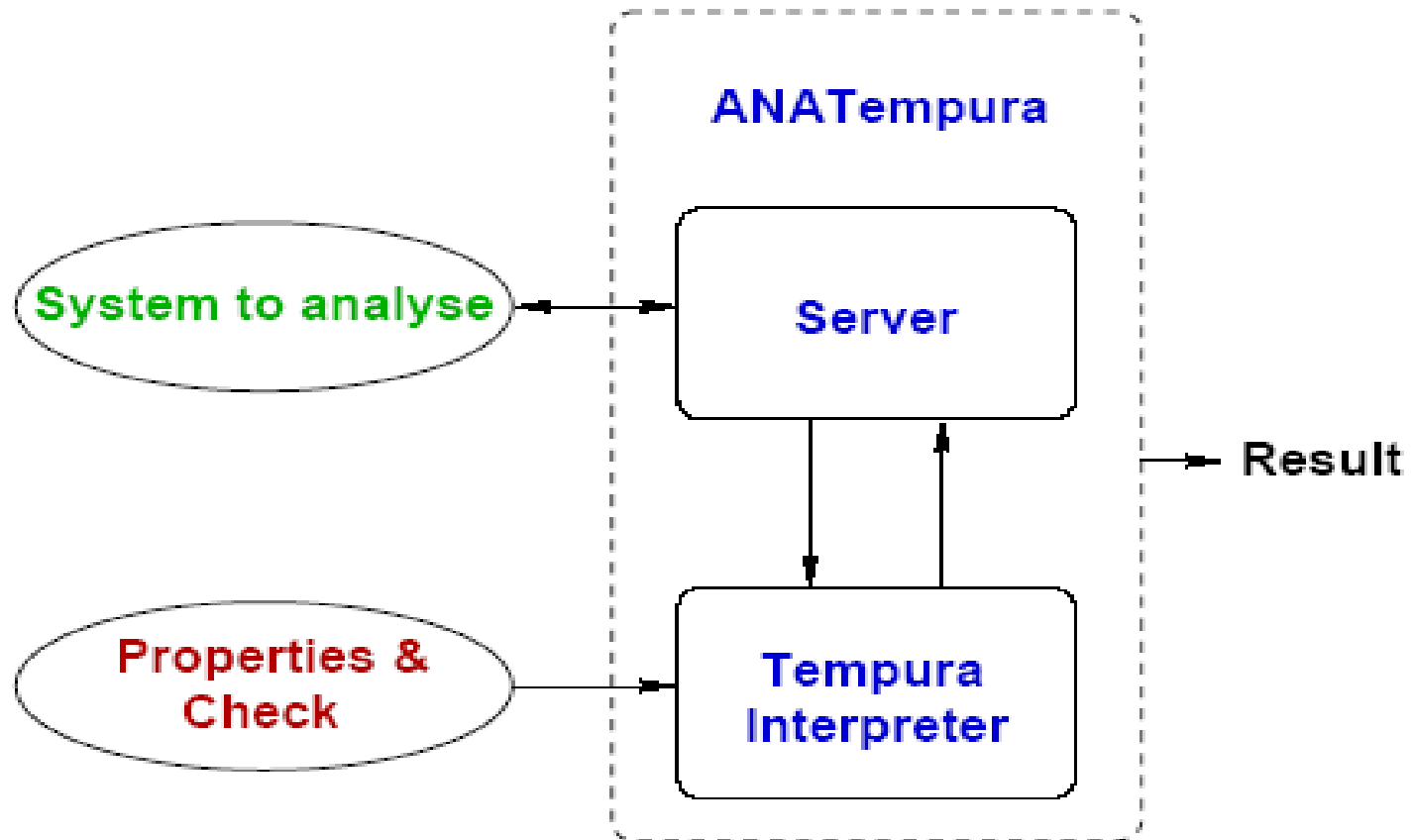
Choosing a proper formalism

- ▶ Interval Temporal Logic (ITL)
 - Ben Moszkowski – Computer Laboratory, University of Cambridge
- ▶ What is ITL:
 - First order logic with added time dependent operators like “sometimes”, “always”, “next” ...
 - Considering time as a discrete sequence of points in time called intervals.
 - For ITL there is an interpreting mechanism and its program realisation called Tempura.

Tempura

- ▶ Imperative programming language which use subset of ITL:
 - First interpreter was written in Prolog
 - C/C++ version:
 - Roger Hale, Ph.D. thesis in Cambridge , 1984–1985 г.,
 - Maintenance: Antonio Cau, STRL, De Montfort University.
- ▶ AnaTempura
 - The centralize surrounding environment of Tempura.

AnaTempura



Approach

- ▶ Three possible ways:
 - Wrapping Tempura with I/O Java classes.
 - Creating a complete new Java version of ITL interpreter.
 - Reengineering the existing C-based version of Tempura.

Why reengineering?

- ▶ Missing documentation and specification of the basic algorithms used in the interpreter.
- ▶ Homogeneous environment.
- ▶ Using proven system and already prepared test cases.

Reengineering in steps

- ▶ Iterative hand-made translation
 - C to Java without changing the imperative structure of the system.
 - Imperative Java to OO Java (jTempura).
 - OO Java to AO Java (JADE based AjTempura) following VES specification.

From jTempura to AjTempura

- ▶ To support the reengineering process we create a model:
 - C3A model
 - Abstract model for Context-Aware Agent Architecture
 - Functionality is supported by agents (persistent agents *PA* and operative agents *OA*).
 - PA supports the basic functionality of the system.
 - OA are generated by PA.

More details : AjTempura – First realization of C3A model, IEEE IS`14, Warsaw, Poland 2014

C3A Lifecycle

repeat

running(a_1), **running**(a_2), ..., **running**(a_p);

anytime $\forall a_i$ {

if ($(\exists e_k \in R(e_k)) \wedge (R(e_k) = R(a_i))$)

then {

$a_j \leftarrow \text{GENERATE}(a_j, e_k)$;

send($a_i, a_j, \text{REQUEST}(e_k)$);

when **INFORM** ($a_j, a_i, \text{'done'}$)

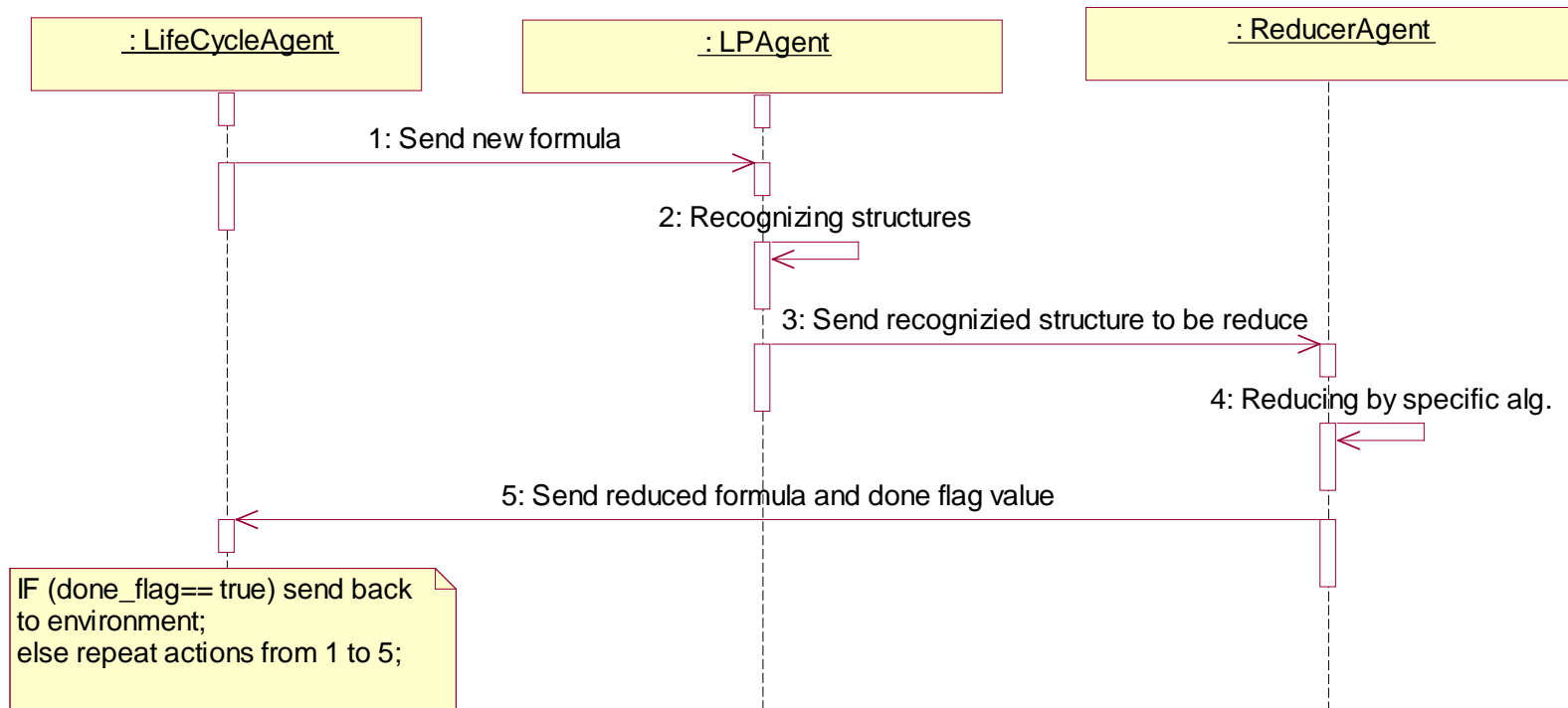
then (**REMOVE**(a_i, a_j) \vee **SELFREMOVE**(a_j)) }

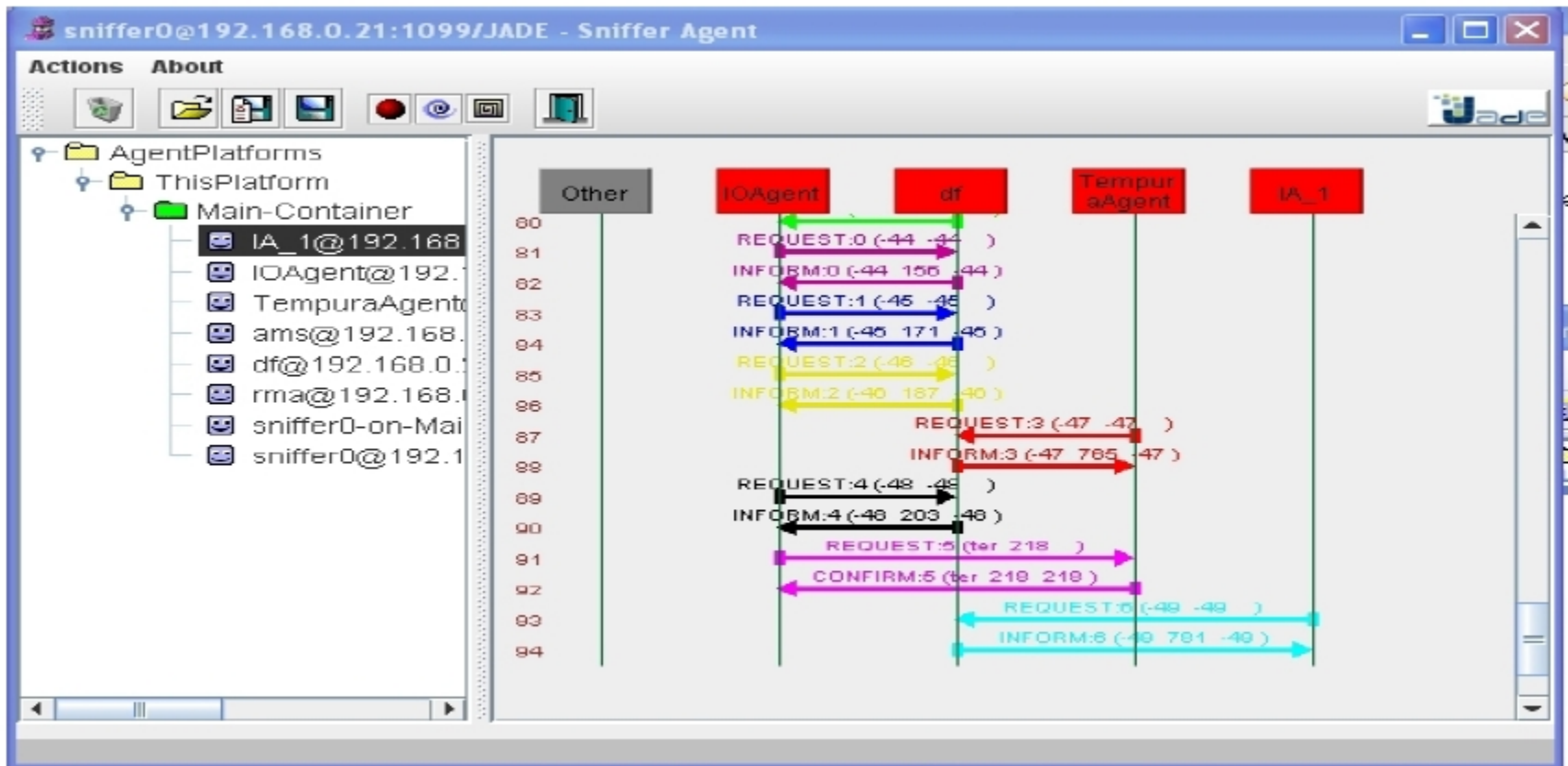
endif

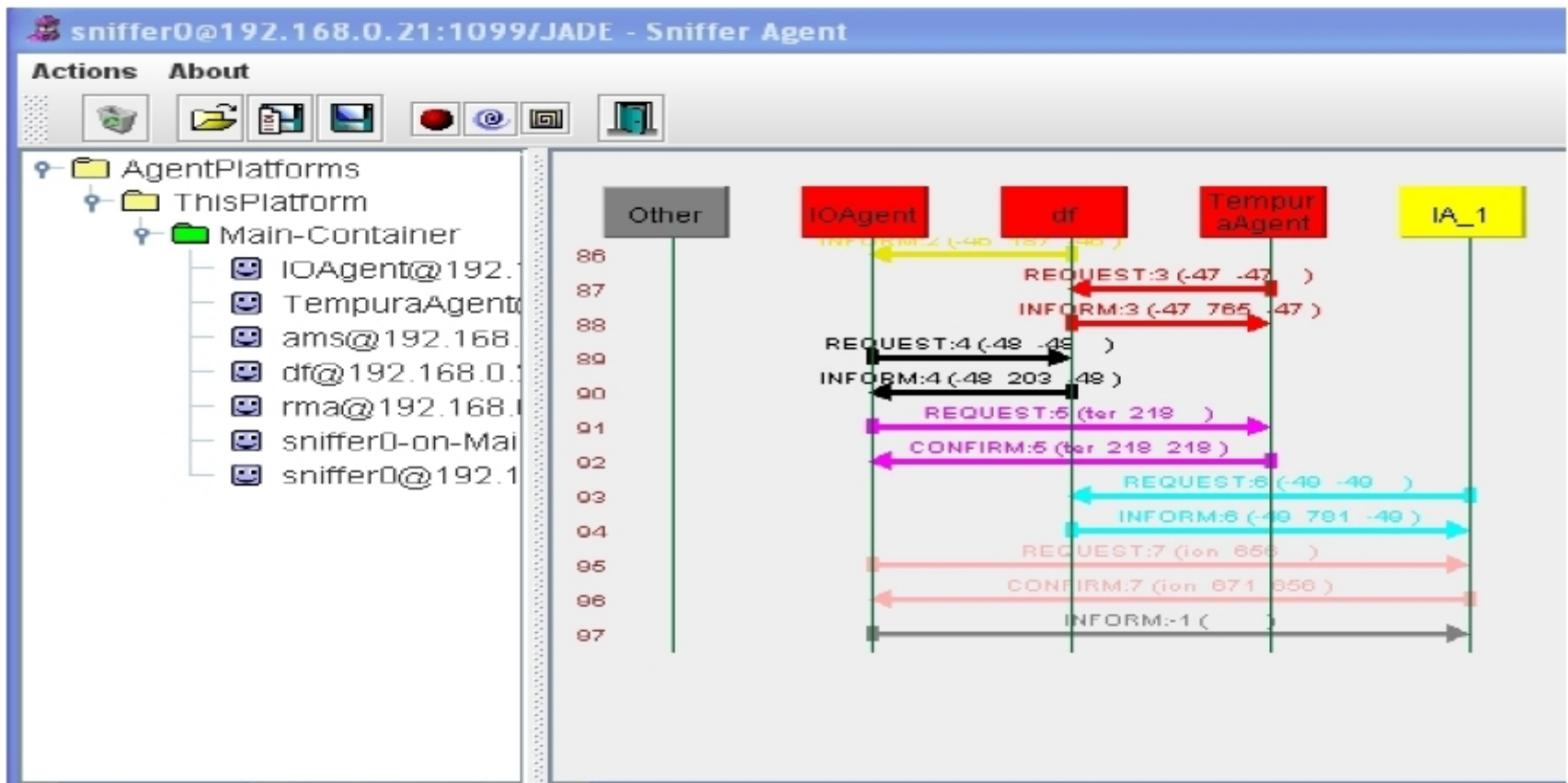
endanytime

forever

AjTempura life-cycle







**THANK YOU FOR THE
ATTENTION**